| Questions: |
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| 1. Consider the foll <br> with start symbol X <br> 1. $\mathrm{X} \rightarrow(\mathrm{P})$ <br> 2. $\mathrm{P} \rightarrow \mathrm{ZP}$ <br> 3. $\mathrm{P} \rightarrow \mathrm{Z}$ <br> 4. $\mathrm{Z} \rightarrow 0$ <br> 5. $\mathrm{Z} \rightarrow 1$ |

a) Attempt to give a table for tabledriven parsing by filling it in with proper entries as much as possible and by placing multiple entries in cells where more than one would be needed.
b) Which terminal symbols are in FIRST(ZP) $\cap \operatorname{FIRST}(Z)$ ?
(Your table should tell you.)
c) Why is the grammar not $\operatorname{LL}(1)$ ?
2. Using the technique of creating tails with $\varepsilon$ :
a) convert the grammar in Problem 1 to an $\operatorname{LL}(1)$ grammar with $\varepsilon$.
b) Give a table for table-driven parsing for your new grammar.

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| 3. Create a grammar that describes a grammar rule for a context free grammar (without $\varepsilon$ ). The syntax of the rule for which you are to give a grammar should be like those in Problem 1 (capital letters for nonterminals, non-capitals for terminals, $\rightarrow$ for "produces," and \| for "or"). As a simplification, assume that the only non-terminals are $\mathrm{A}, \mathrm{B}$, and C , and that, besides the meta-symbols $\rightarrow$ and \|, the only terminal symbols are $\mathrm{x}, \mathrm{y}$, and z . <br> Comments: Observe that because you are giving a grammar for a grammar rule, you will have problems with the meta-symbols. (For example, the meta-symbol \| can't be both a terminal symbol and a separator for alternative right hand sides.) Resolve this problem as follows. Use BNF syntax for your grammar. Thus, the meta-symbol $\rightarrow$ can be a terminal since it is not a meta-symbol in the BNF syntax. Further, you can use a non-terminal such as <bar> to denote the meta-symbol \|. Add the rule" <bar> ::= \|", which cannot be ambiguous, to give <bar> its standard terminal symbol. Even though the grammar rule you are describing has no $\varepsilon$, the grammar that describes the grammar rule may have $\varepsilon$. |  |
| 4. Using the statements R and H respectively for "Mark is rich" and "Mark is happy," write the following statements in symbolic form. <br> a) Mark is not rich. <br> b) Mark is rich and happy. <br> c) Mark is rich or happy. <br> d) If Mark is rich, then he is happy. |  |


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| 5. Identify the atomic propositions of <br> the following sentences and <br> replace them by propositional <br> symbols. Then translate the <br> sentences into propositional <br> calculus. |  |
| a) If you do not leave, I will call the |  |
| police. |  |
| b) I am sad if and only if I am not |  |
| happy. |  |
| c) It is a nice day if it is sunny and it |  |
| is not hot. |  |
| d) if $i>j$, then $i-1>$ j, else (if i is |  |
| not $\gg j) j=3$ |  |$|$



